

ECE 356/COMPSI 356

Computer Network Architecture

Introduction to Internetworking.
Switching technologies.

Wednesday September 18, 2019

Recap

- Previous lectures:
 - Ethernet, multiple access
 - Wireless communications
- Readings for this class: **PD 3.1**

Lecture Outline

- Internetworking: an introduction
- Switching
- Types of switching
 - Datagram switching
 - Virtual circuit switching
 - ATM
 - Source routing

Internetworking

- Previously we saw how to connect one node to another, or to an existing network. How do we build networks of global scale?
- How do we interconnect different types of networks to build a large global network?
- The foundation of the Internet

Next 6 Lectures: A Roadmap

1. Switching and bridging
2. Ethernet switches
3. Basic Internetworking (IP)
4. Routing – Distance Vector-based
5. Routing – Link State-based

Lecture Outline

- Internetworking: an introduction
- **Switching**
- Types of switching
 - Datagram switching
 - Virtual circuit switching
 - ATM
 - Source routing

Why Do We Need Switches?



- Problem: single link networks have limited scale
 - Ethernet < 1024 hosts, 2500 meters
 - Wireless limited by radio ranges
 - Point-to-point links connect only two nodes

Packet Switching

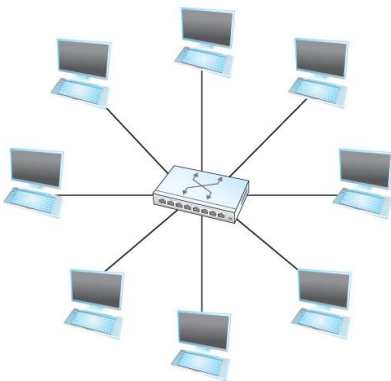
- A **packet switch** is a device with several inputs and outputs leading to and from the nodes that the switch interconnects
 - Hosts communicate without being directly connected
- Main responsibility: receive incoming packets on one of its links and to transmit them on some other link

Difference Between a Switch and a Hub

- Video posted on Piazza

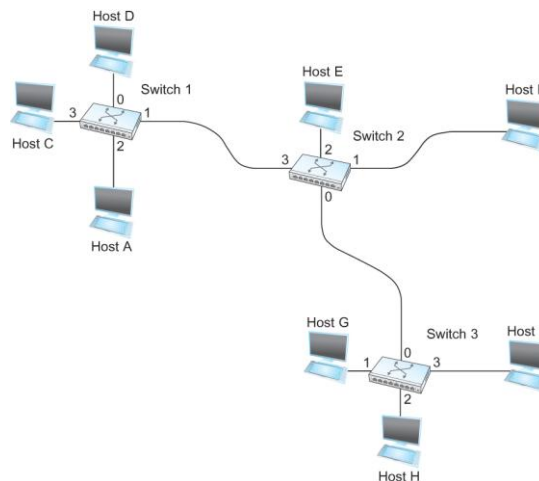
9

A Star Topology



- A switch has a limited number of input and output ports
- Adding a new host may not reduce the performance for other hosts
 - Not true for shared media networks
 - Why?

Switches Can Be Connected to Each Other to Build Larger Networks



Duke UNIVERSITY

Switching Technologies

- Switching / forwarding: receiving incoming packets on one of switch links and transmitting them on some other link
- Problem: how does a switch decide on which output port to place each packet?
- Solution: looks at the packet header and makes a decision
 - Connectionless: datagram
 - Connection oriented: virtual circuit
 - Source routing

Duke UNIVERSITY

Challenges

- Contention
 - Input rate exceeds output rate
 - Multiple input ports may send to the same output port
 - Switches queue packets until contention disappears
- Congestion
 - When a switch runs out of buffer, it discards packets
 - Too frequent packet loss is said to be congested

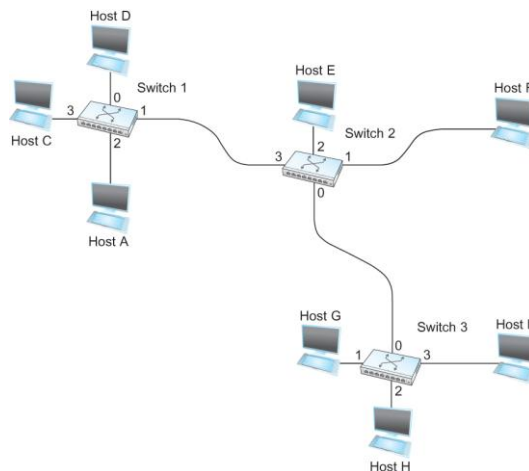
Lecture Outline

- Internetworking: an introduction
- Switching
- Types of switching
 - **Datagram switching**
 - Virtual circuit switching
 - ATM
 - Source routing

Datagram

- Every packet contains the complete destination address
 - A global unique identifier
 - Ethernet has 48-bit addresses
- A switch maintains a **forwarding table** that maps a packet to an output port

Example: Switch 2 Forwarding Table



| | |
|---|---|
| A | 3 |
| B | 0 |
| C | 3 |
| D | |
| E | |
| F | |
| G | |
| H | |

Routing is the process of creating forwarding tables

Features of Datagram Switching (1/2)

- Connectionless
 - Hosts can send anytime. No need to wait for connection to set up
- Unknown network state
 - Not sure whether a packet can reach the destination
- Independent forwarding
 - Packets can take different paths

Features of Datagram Switching (2/2)

- Robust to failures
 - A failure of a switch may not disrupt communications
 - Switches can re-compute forwarding tables
- Commonly deployed

Lecture Outline

- Internetworking: an introduction
- Switching
- Types of switching
 - Datagram switching
 - **Virtual circuit switching**
 - ATM
 - Source routing

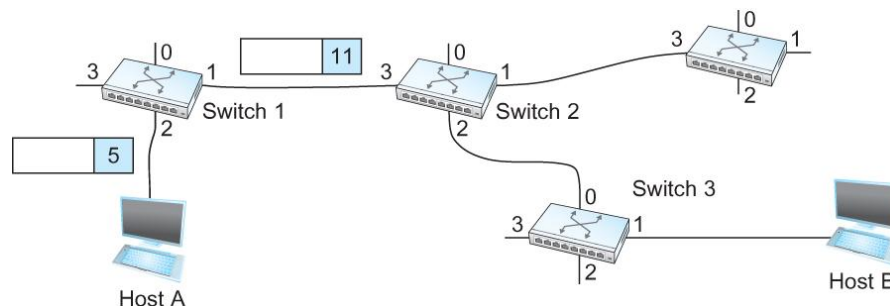
Virtual Circuit Switching

- *Connection-oriented* model
- Two steps:
 - Set up a virtual circuit
 - Data transfer

Virtual Circuit Switching: Connection Setup Phase

- Set up *connection state* in each of the switches between the source and the destination
- **VC table** entry in each switch, for each VC:
 - A virtual circuit identifier (VCI)
 - An incoming interface
 - An outgoing interface
 - An outgoing VCI

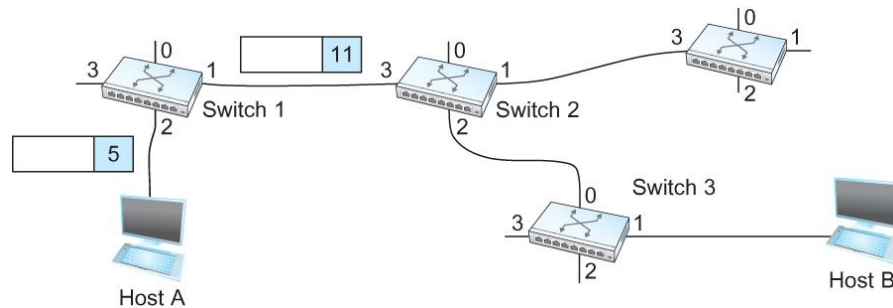
Virtual Circuit Switching: Algorithm



- If a packet arrives on the matching incoming port with the matching incoming VCI, it will be sent to the corresponding outgoing port with the corresponding VCI
- VCIs are link-local → VCI IDs can be kept short

Example: Virtual Circuit Table, Switch 1

| Incoming interface | Incoming VCI | Outgoing interface | Outgoing VCI |
|--------------------|--------------|--------------------|--------------|
| 2 | 5 | 1 | 11 |



Duke UNIVERSITY

How to Establish Connection State

- Administrator-configured
 - **Permanent virtual circuit (PVC)**
 - Administrators manually sets up VC tables
 - Does not suit large networks
- Signaling to establish state
 - **Switched virtual circuit (SVC)**
 - A host sends messages to dynamically setup or tear down a VC

Duke UNIVERSITY

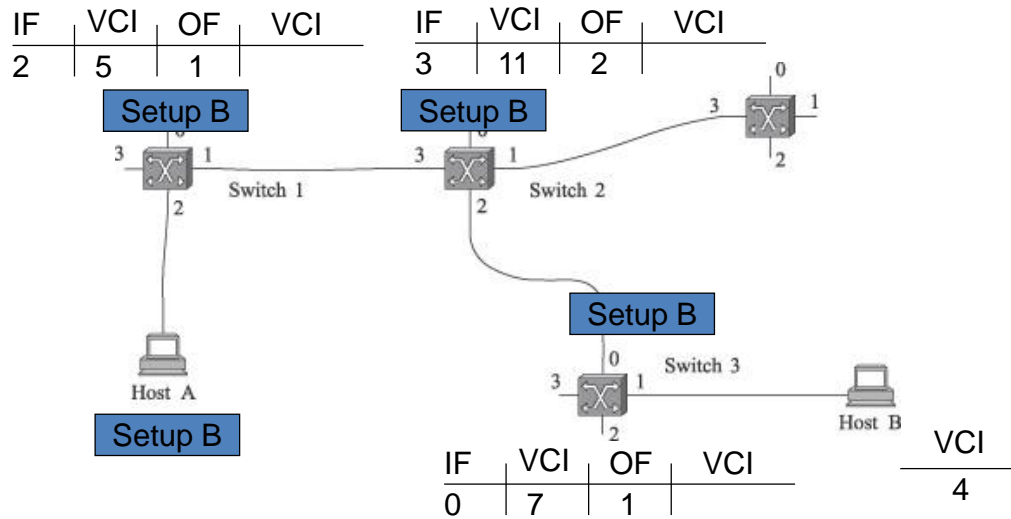
VC Setup Protocol: Host A to Destination B (1/2)

- A host A sends a setup message to first hop switch, including the final destination address
 - Similar to a datagram packet
 - Message to get all the way to B
 - Assume switches know how to do it
- The switch picks an unused VCI to identify the incoming connection, and fills part of the VC table
 - *Why not let the host pick the VCI value?*

VC Setup Protocol: Host A to Destination B (2/2)

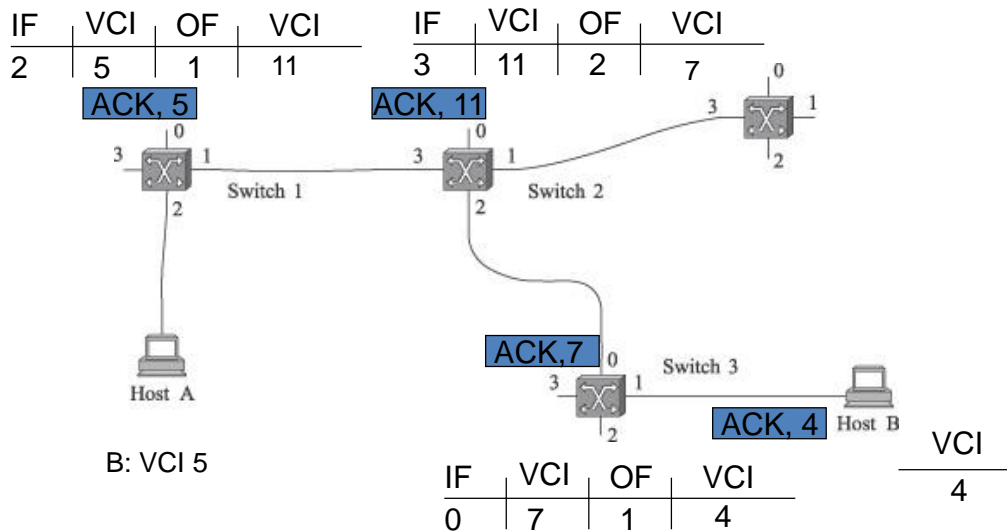
- Every switch repeats the process until the packet reaches the destination B
- The destination B sends an ACK to inform its upstream switch the VCI for the connection

VC Setup: An Example (1/2)



Duke UNIVERSITY

VC Setup: An Example (2/2)



Duke UNIVERSITY

After VC Setup is Complete

- Data transmission can begin
- Unused connections need to be torn down

29

Characteristics of VC Switching

- Pros:
 - Data packets contain a small VCI, not the full destination addresses
- Cons:
 - Connection setup wait
 - One switch failure tears down the entire connection
 - Connection sets up require routing algorithms
 - Setup packet is forwarded using a datagram algorithm

VC Allows Resource Reservation

- Pro: buffers can be allocated during the setup phase to avoid congestion
- An example (X.25)
 - Buffers allocated during connection setup
 - Sliding window is run between pairs of nodes: *hop-by-hop flow control*
 - Circuit is rejected if no more buffer

Quality of Service (QoS)

- Connectionless network is difficult to allocate resources
 - Switches send packets independently
 - How to associate one packet with other packets?
- Virtual circuit can be used to provide different QoS
 - Allocate a fraction of link bandwidth to each circuit
- Examine QoS in detail in the second half of the course

Technologies That Use VC

- A number have been popular over the years
 - X.25 – financial transaction systems (ATMs, point-of-sale terminals, credit card payments), aviation
 - Frame Relay
 - Asynchronous Transfer Mode (ATM)

Lecture Outline

- Internetworking: an introduction
- Switching
- Types of switching
 - Datagram switching
 - Virtual circuit switching
 - **ATM**
 - Source routing

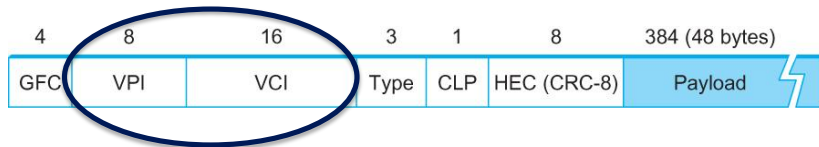
Asynchronous Transfer Mode

- *ATM Cells*: fixed-size packets
 - 5 bytes header
 - 48 bytes payload
- If payload smaller than 48B, uses padding
- If greater than 48B, breaks it

Why Small, Fixed-Length Packets?

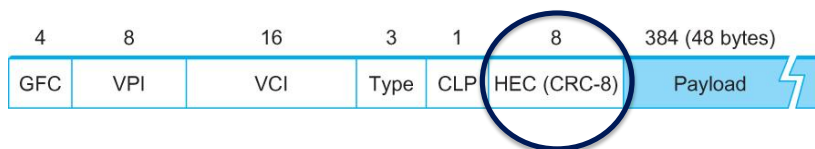
- Cons: maximum efficiency $48/53=90.6\%$
- Pros:
 - Suitable for efficient high-speed hardware implementation
 - Many switching elements doing the same thing in parallel

ATM: Virtual Paths



- 24-bit circuit identifiers
- Two-levels of hierarchy
 - 8-bit Virtual Path Identifier (VPI), 16-bit VCI
 - Virtual paths shared by multiple connections
 - Circuits with the same VPI could be switched together

ATM: Header Error Check



- 8-bit CRC specifically for the header
 - Error causes an ATM cell to be mis-delivered

History of ATM: Why 48 Bytes?

- It's from the telephone technology
- Thought data would be mostly voice
- A compromise
 - US wanted 64 bytes for efficiency
 - Europe wanted 32 bytes for simplifying echo cancellation
 - $(64+32) / 2 = 48$ bytes

History of ATM: Where Is It Now?

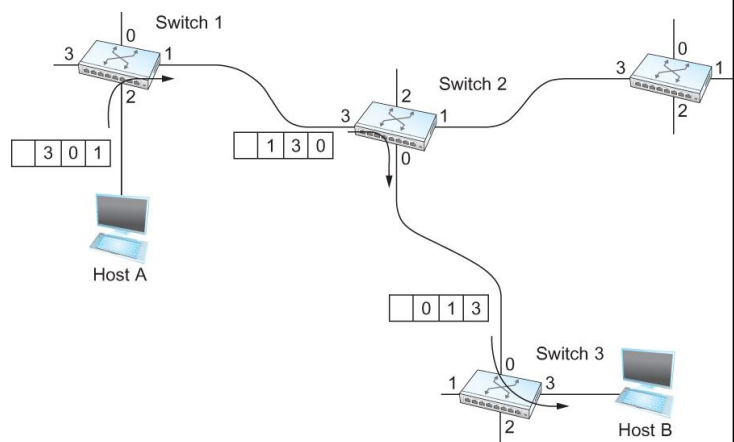
- Popular in the late 80s and early 90s due to its high speed
 - Major telecommunication companies supported it
- Popularity faded. IP/Ethernet ruled
 - IP over ATM
 - DSL over ATM: DSL modem takes Ethernet frames and chop them into cells

Lecture Outline

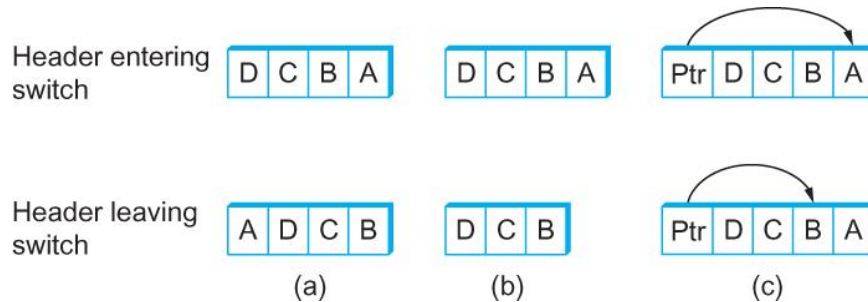
- Internetworking: an introduction
- Switching
- Types of switching
 - Datagram switching
 - Virtual circuit switching
 - ATM
 - **Source routing**

Source Routing

- Source host provides all the information for packets to travel across the network
 - Packets carry output port numbers
 - Packets carry switch addresses
 - Variable header length



Handling Source Routing Headers



- a. Rotation
- b. Stripping
 - No return path!
- c. Pointer

Loose or Strict Source Routing

- Strict
 - Must visit every node on the path
- Loose
 - Waypoints rather than the complete route
 - Impractical to expect to be able to construct the entire route in a large network

Summary

- Internetworking: an introduction
- Switching
- Types of switching
 - Datagram switching
 - Virtual circuit switching
 - ATM
 - Source routing

Next Lecture

- Switches that are used to forward packets between local area networks
 - LAN switches
 - Bridges